

## Priority Queues

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## Outline and Reading

- ◆ PriorityQueue ADT (§7.1)
- ◆ Total order relation (§7.1.1)
- ◆ Comparator ADT (§7.1.4)
- ◆ Sorting with a priority queue (§7.1.2)
- ◆ Selection-sort (§7.2.3)
- ◆ Insertion-sort (§7.2.3)

## Priority Queue ADT

- ◆ A priority queue stores a collection of items
- ◆ An item is a pair (key, element)
- ◆ Main methods of the Priority Queue ADT
  - `insertItem(k, o)` inserts an item with key *k* and element *o*
  - `removeMin()` removes the item with the smallest key
- ◆ Additional methods
  - `minKey(k, o)` returns, but does not remove, the smallest key of an item
  - `minElement()` returns, but does not remove, the element of an item with smallest key
  - `size()`, `isEmpty()`
- ◆ Applications:
  - Standby flyers
  - Auctions
  - Stock market

## Total Order Relation

- ◆ Keys in a priority queue can be arbitrary objects on which an order is defined
- ◆ Two distinct items in a priority queue can have the same key
- ◆ Mathematical concept of total order relation  $\leq$ 
  - Reflexive property:  $x \leq x$
  - Antisymmetric property:  $x \leq y \wedge y \leq x \Rightarrow x = y$
  - Transitive property:  $x \leq y \wedge y \leq z \Rightarrow x \leq z$

## Comparator ADT



- ◆ A *comparator* encapsulates the action of comparing two objects according to a given total order relation
- ◆ A generic priority queue uses a comparator as a template argument, to define the comparison function (`<`, `=`, `>`)
- ◆ The comparator is external to the keys being compared. Thus, the same objects can be sorted in different ways by using different comparators.
- ◆ When the priority queue needs to compare two keys, it uses its comparator

## Using Comparators in C++



- ◆ A comparator class overloads the `()` operator with a comparison function.
- ◆ Example: Compare two points in the plane lexicographically.
- ◆ To use the comparator, define an object of this type, and invoke it using its `()` operator.
- ◆ Example of usage:

```

class LexCompare {
public:
    int operator()(Point a, Point b) {
        if (a.x < b.x) return -1
        else if (a.x > b.x) return +1
        else if (a.y < b.y) return -1
        else if (a.y > b.y) return +1
        else return 0;
    }
};

Point p(2.3, 4.5);
Point q(1.7, 7.3);
LexCompare lexCompare;

if (lexCompare(p, q) < 0)
    cout << "p less than q";
else if (lexCompare(p, q) == 0)
    cout << "p equals q";
else if (lexCompare(p, q) > 0)
    cout << "p greater than q";
    
```

## Sorting with a Priority Queue

- ◆ We can use a priority queue to sort a set of comparable elements
  1. Insert the elements one by one with a series of `insertItem(e, e)` operations
  2. Remove the elements in sorted order with a series of `removeMin()` operations
- ◆ The running time of this sorting method depends on the priority queue implementation

```

Algorithm PQ-Sort(S, C)
Input sequence S, comparator C
        for the elements of S
Output sequence S sorted in
        increasing order according to C
P ← priority queue with
        comparator C
while !S.isEmpty()
    e ← S.remove(S.first())
    P.insertItem(e, e)
while !P.isEmpty()
    e ← P.minElement()
    P.removeMin()
    S.insertLast(e)
    
```

## Sequence-based Priority Queue

- ◆ Implementation with an unsorted sequence
  - Store the items of the priority queue in a list-based sequence, in arbitrary order
- ◆ Performance:
  - `insertItem` takes  $O(1)$  time since we can insert the item at the beginning or end of the sequence
  - `removeMin`, `minKey` and `minElement` take  $O(n)$  time since we have to traverse the entire sequence to find the smallest key
- ◆ Implementation with a sorted sequence
  - Store the items of the priority queue in a sequence, sorted by key
- ◆ Performance:
  - `insertItem` takes  $O(n)$  time since we have to find the place where to insert the item
  - `removeMin`, `minKey` and `minElement` take  $O(1)$  time since the smallest key is at the beginning of the sequence

## Selection-Sort

- ◆ Selection-sort is the variation of PQ-sort where the priority queue is implemented with an unsorted sequence
- ◆ Running time of Selection-sort:
  1. Inserting the elements into the priority queue with  $n$  `insertItem` operations takes  $O(n)$  time
  2. Removing the elements in sorted order from the priority queue with  $n$  `removeMin` operations takes time proportional to  $1 + 2 + \dots + n$
- ◆ Selection-sort runs in  $O(n^2)$  time

## Insertion-Sort

- ◆ Insertion-sort is the variation of PQ-sort where the priority queue is implemented with a sorted sequence
- ◆ Running time of Insertion-sort:
  1. Inserting the elements into the priority queue with  $n$  `insertItem` operations takes time proportional to  $1 + 2 + \dots + n$
  2. Removing the elements in sorted order from the priority queue with a series of  $n$  `removeMin` operations takes  $O(n)$  time
- ◆ Insertion-sort runs in  $O(n^2)$  time

## In-place Insertion-sort

- ◆ Instead of using an external data structure, we can implement selection-sort and insertion-sort in-place
- ◆ A portion of the input sequence itself serves as the priority queue
- ◆ For in-place insertion-sort
  - We keep sorted the initial portion of the sequence
  - We can use `swapElements` instead of modifying the sequence

